

Preserve the Privacy of Anonymous and Confidential DB using K-anonymity

Vishakha S. Kulkarni, Sonali Patil

Department of Computer Science, Alard College of Engineering & Management, Pune, Maharashtra, India

ABSTRACT

In IT field, maintaining confidentiality and privacy of data is very important for decision making. And there is need of exchanging certain data and published with respect to the demand. The information that is to be exchanged contains private and sensitive information which moves between various parties and this may violate the privacy of data. Thus, privacy preservation of information in its accurate form while moving between various parties, my goal is providing mechanism that will not allow the unauthenticated user to modify the data known as k_anonymous technique. Two protocols are used in this application that will solve this problem based on two methods suppression and generalization. Some cryptographic assumptions are made to use these protocols, and theoretical analysis is provided to prove their experimental results that illustrate their efficiency.

Keywords: Anonymity, data management, privacy, secure computation.

I. INTRODUCTION

The database contains important information and is valuable asset for many applications. Thus their security is important. Data confidentiality is relevant because of the value that data have. The medical data of patients collected over several years represent history of patients and is a valuable data that needs to be protected. This requirement gave rise to a large variety of approaches that aim at protecting data confidentiality and data ownership. Data confidentiality problems are created by an unauthorized user to get the knowledge about data stored in the database. Access to individual's personal information is limited by privacy. It means the authorized access by valid users.

Database privacy follows confidentiality, availability, and integrity of personal data and not confidentiality alone. Anonymization is the technique that is required to provide privacy. Anonymization is the method of masking data. In this from original data, identifying information is removed so that to protect personal or private information. Using data Anonymization information can be transferred between two organizations, by converting data in to cipher data using encryption method. K-Anonymization is one of the

approaches that maintain privacy of data. In K-Anonymization approach, at least K-tuples should be indistinguishable by masking values.

The data providers are medical facilities or college that provide sensitive information related to patient or staff through anonymous authentication and connection. User is authenticated done using user ID and password. The data provider's data privacy is to be protected from outside valid users like researcher's or staff as the database is in anonymous form.

The existing system face some difficulties concerning that the without revealing the contents of tuples and database by users, how data integrity can be preserved by developing the anonymity of DB. Algorithms are established to provide database anonymization. It relates with how data privacy of whole databases and their owner and also individual tuples and its owner is maintained without disclosing the contents.

This system has, suppression based as well as generalization based anonymous database. A secure protocol is presented that privately checks whether K-anonymous database retains its anonymity even after insertion of a new tuple.

II. METHODS & MATERIALS

The information concerning a data provider is stored in a single tuple, and DB is kept confidentially at the server. Since DB is anonymous, the data provider's privacy is protected from researchers. Such task is guaranteed through the use of anonymization. Preserving the privacy & confidentiality without revealing the contents of tuple and DB is done by establishing the anonymity of DB. A secure protocol is presented for privately checking whether K-anonymous database remains anonymous even after insertion of a new tuple. Suppressed the value of attribute by replacing "*" and Generalized the value with related possible general value to maintain the k-anonymity in database. Thus by making such k-anonymity in table it becomes complicated for third party to identify the record. In the system, before a tuple is inserted the data can be encrypted using shared secret key AES algorithm. Based on a commutative encryption function the data provider can share a secret key with each other using Diffie-Hellman Algorithm.

As shown in Fig. 7.1, proposed system consists of following modules:

- Admin Module.
- Data Provider for Suppression and Generalization.
- Server for Suppression and Generalization.

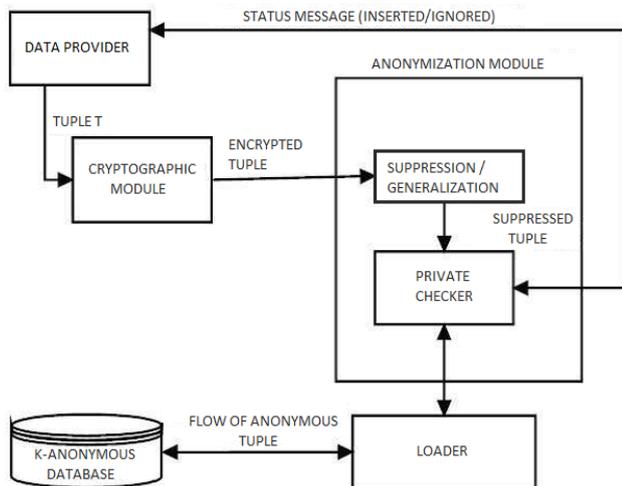


Figure 1: Proposed System Architecture

In this proposed model a secure protocol is presented that privately checks whether database remains k-anonymous even after insertion of new tuple. Quasi-Identifier (QI): QI is a minimal set of attributes which is used to uniquely identify individuals. Attack is mainly using Quasi-Identifier. Attacks may be re-identification

or linking attack. To prevent the attack, masks the values of Quasi-Identifiers using either suppression based or Generalization based Anonymization methods. In Suppression based anonymization method, mask the Quasi-Identifiers value using a special symbol like * and in Generalization based anonymization method, replace a specific value with a more general one using Value Generalization Hierarchies (VGH).

The diffiehellman key exchange algorithm is used to generate private secure key. Then AES algorithm is applied to encrypt and decrypt data by using the key generated by the diffiehellman key exchange algorithm. When user enters his information then this information is encrypted by using AES and also all data in table is encrypted using same algorithm. If information from user matches with table information the tuple will be decrypted and inserted into table.

Let the data provider is X and Suppression & Checker module is Y. The flow of operation is given below:

- a. X sends a tuple T in to cryptographic module.
- b. The cryptographic module encrypts the tuple (encryption means to convert the plane text in to ciphertext) and send it to the Suppression & Checker module(Y)
- c. Y, then compute.
- d. The anonymized version of tuple T.
- e. Check whether the data is matched with data's in the loader.
- f. The loader reads chunks of anonymized tuple from the K-Anonymous database.
- g. If the tuples are not matched, then the loader reads next chunks of anonymized tuple from the k-anonymous database and checking can be performed.
- h. If any match found, then the tuple t can be inserted in to the K-Anonymous database.
- i. Finally we can send a message to the data provider about the status of the tuple T (status are INSERTED/ IGNORE).
- j. According to the status, the data provider can decide further action.

A. Module1: Admin Module:

Module 1 is the Admin module. Admin enters username and password; if it is correct then admin is validated and can view original data and anonymous data report further. If admin enters wrong information then user is invalid user and can't proceed further.

**B. Module2: Data Provider for Suppression Method
Data Provider for Generalization Method**

Anonymization is technique which hides sensitive attribute value in such a way that it cannot be identified back. In k-anonymization approach the total number of rows is k and k cannot be differentiated with other k-1 rows by taking into account only a set of attributes, then this table is known as K-anonymized. Privacy preservation can be done by simply using k-anonymization approach on suppression and generalized techniques.

a) Suppression Based Anonymous Technique:

When suppression-based anonymization method is used, consider a table $T = \{t_1, \dots, \dots\}$ tuples over the attribute set A. In suppression method, the values of some well-chosen attributes are masked to form subsets. It is mask with the special value „*“. Forming the subset and classify that subsets by using Quasi-Identifier (QI). Quasi-Identifier (QI): Each record contains a number of attributes: some attributes are unique and personal attributes (such as disease and salary) and some may be repeated and general that is quasi-identifiers (called QI, such as zipcode, age, and gender) by taking this it can easily identify someone. Consider the example of patient As shown in table 1 which contains original database (Table T) having Quasi-Identifier $QI = \{Zipcode, age, Nationality\}$ or more sensitive three attributes value. After applying suppression based technique on original dataset the original dataset is anonymized and Table 2 shows a suppression based k-anonymization with $k=2$ it means that at least $k=2$ tuples should be indistinguishable by masking values.

b) Generalization Based Anonymous Technique

In generalization-based anonymization consists in substituting the values of a given attribute with more general values in the database, according to a priori established value generalization hierarchies (VGHs) with some Cryptographic Primitives. In Table 1 original information is stored and after performing generalization techniques on original dataset the original dataset is anonymized and table 2 gives generalized data with $k=3$. The Generalization is technique which replaces a value with a “less-specific but semantically consistent” value. It can be defined based on the VGH which specify how

the data will be generalized. According to the VGH of DISEASE, say that the value of disease is generalized according to the disease causes. Like “HIV” cause by virus so it can be generalized to “Diseases Caused by virus”. The attribute Age is generalized to the interval (30-39).

Table -1: Original Patient Data

	<i>Zip code</i>	<i>Age</i>	<i>Nationality</i>	<i>Condition</i>
1	13053	28	Russian	Heart disease
2	13068	29	American	Heart disease
3	13068	21	Japanese	Viral infection
4	13053	23	American	Viral infection
5	14853	50	Indian	Cancer
6	14853	55	Russian	Heart disease
7	14850	47	American	Viral infection
8	14850	49	American	Viral infection
9	13053	31	American	Cancer
10	13053	37	Indian	Cancer
11	13068	36	Japanese	Cancer
12	13068	35	American	cancer

Table -2: Anonymous Patient Data

	<i>Zip code</i>	<i>Age</i>	<i>Nationality</i>	<i>Condition</i>
1	130**	<30	*	Heart disease
2	130**	<30	*	Heart disease
3	130**	<30	*	Viral infection
4	130**	<30	*	Viral infection
5	1485*	≥40	*	Cancer
6	1485*	≥40	*	Heart disease
7	1485*	≥40	*	Viral infection
8	1485*	≥40	*	Viral infection
9	130**	3*	*	Cancer
10	130**	3*	*	Cancer
11	130**	3*	*	Cancer
12	130**	3*	*	Cancer

**C. Module3: Server for Suppression Method
Server for Generalization Method**

In this module, the suppressed tuple is compared by the tuple loaded from k-anonymous database in loader. Private checker compares this both the tuple, if they are same then the tuple is inserted. Otherwise the tuple is ignored. The system actually updates the database depends on the result of the anonymity checker. In some cases the insertion or updation failed in k-anonymous

database then it waits until k-1 value becomes positive and other tuples fail the insertion.

III. RESULTS and DISCUSSION

A Private Checker is composed by the following modules: a crypto module that is in charge of encrypting all the tuples exchanged between a user and the Private Updater, a checker module that performs all the controls, a loader module that reads chunks of anonymized tuples from the k-anonymous DB. The chunk size is fixed in order to minimize the network overload. The functionality provided by the Private Checker prototype regards the check on whether the tuple insertion into the k-anonymous DB is possible. The information flow across the above mentioned modules is as follows: after an initial setup phase in which the user and the Private Checker prototype exchange public values for correctly performing the subsequent cryptographic operations, the user sends the encryption of her/his tuple to the Private Checker; the loader module reads from the k-anonymous DB the first chunk of tuples to be checked with encrypted tuple. Such tuples are then encrypted by the crypto module. The checker module performs the above mentioned check one tuple at time in collaboration with the user. If none of the tuples in the chunk matches the User tuple, then the loader reads another chunk of tuples from the k-anonymous DB.

IV. CONCLUSION

Data confidentiality and privacy is a challenging problem faced in case of security of database. In this work, two secure protocols are presented for privately checking whether a k-anonymous database retains its anonymity once a new tuple is being inserted to it. Since the proposed protocols ensure the updated database remains k-anonymous. The data provider's privacy cannot be violated from any user's updating the table. So the database is updated properly using the proposed protocols. This is useful in medical application. If insertion of record satisfies the k-anonymity then such record is inserted in table and suppressed the sensitive information attribute by "*" to maintain the k-anonymity in database. Thus, by making such k-anonymity in table that makes unauthorized user too difficult to identify the record.

V. REFERENCES

- [1]. Sivasubramanian .R, K.P. Kaliyamurthie, "Privacy-Preserving Updates to Anonymous Databases", IJCSMC, Vol. 2, Issue. 4, April 2013, pg.582 – 587.
- [2]. Alberto Trombetta, Wei Jiang, Member, IEEE, Elisa Bertino, Fellow, IEEE, and Lorenzo Bossi, "Privacy-Preserving Updates to Anonymous and Confidential Databases" IEEE Transactions on Dependable and Secure Computing, VOL. 8, No. 4, July/August 2011.
- [3]. L. Sweeney, "K-Anonymity: A Model for Protecting Privacy", International Journal on Uncertainty, Fuzziness and Knowledge-based Systems, 10 (5), 2002.
- [4]. Dr. K. P.Thooyamani, Dr.V.khanaa, Er.M.R.Arun Venkatesh, "Privacy-Preserving Updates to Anonymous and Confidential Database", International Journal of Data Mining Techniques and Applications, Volume: 01 Issue: 01 January-June 2012.
- [5]. Ishwarya M. V, Dr. Ramesh Kumar. K, "Privacy Preserving Updates for Anonymous and Confidential Databases Using RSA Algorithm", International Journal of Modern Engineering Research (IJMER), Vol.2, Issue.5, Sep.-Oct. 2012.
- [6]. Yehuda Lindell, Benny Pinkas, "Secure Multiparty Computation for Privacy-Preserving Data Mining", May 6, 2008.
- [7]. Rajeshwari Suryawanshi, Sulabha Patil,"Privacy Preserving updates to Personalized Anonymity Based Anonymous and Confidential Database", International Journal of Advanced Research in Computer Science and Software Engineering, Volume 3, Issue 10, October 2013.
- [8]. Rakesh Agrawal, Alexandre Evfimievski, Ramakrishnan Srikant, "Information haring Across Private Databases", IBM Almaden Research Center June 9, 2003.
- [9]. Mahendrababu P, Rajarajan G,"Apprising in Secured Manner to Anonymous and Confidential Databases", International Journal of Engineering Research & Technology (IJERT), Vol. 2 Issue 1, January- 2013.
- [10].Ebin P.M, Brilley Batley. C, "Privacy Preserving Suppression Algorithm for Anonymous Databases", International Journal of Science and Research (IJSR), Volume 2 Issue 1, January 2013.
- [11].Mr. Mahesh T.Dhande, Mrs. Neeta A.Nemade, "Performance Improvement of Privacy Preserving in K-anonymous Databases Using Advanced Encryption Standard Technique", International Journal of Advanced Research in Computer Science and Software Engineering, Volume 3, Issue 6, June 2013.
- [12].Rajeshwari Suryawanshi, Prof.Parul Bhanarkar, Rashtrasant Tukdoji Maharaj, "Survey On Privacy Preserving Updates On Anonymous Database", International Journal of Engineering Research & Technology (IJERT),Vol. 2 Issue 1, January- 2013.
- [13].Neha Gosai, S.H.Patil, "Generalization Based Approach to Confidential Database Updates", International Journal of Engineering Research and Applications, Vol. 2, Issue 3, May-Jun 2012.
- [14].Lavanya.Gunasekaran, R. Sujatha,"Enhanced Privacy Preserving Updates for Anonymous and Confidential Databases", International Journal of Computer Networks and Wireless Communications (IJCNWC), Vol.2, No.2, April 2012.
- [15].Deepa.B, Meena.R, "Privacy Control Methods for Updating Confidential Databases", International Journal of Soft Computing and Engineering (IJSCE), Volume-1, Issue-ETIC2011, January 2012.