A Review: Area and Delay Efficient Pre-encoded multipliers Based on Non-Redundant Radix-4 Encoding

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ABSTRACT

During this paper, we tend to introduce an efficient design of pre-encoded multiplier. The radix-4 standard multiplier will be accustomed implement quick PC applications, e.g., RSA cryptosystem and to scale back the quantity of iterations and pipelining. The performance of the present algorithms is primarily determined by the economical implementation of the standard multiplication and exponentiation. Mentioned a Booth’s Radix-2 multiplier and calculated its delay, space and power. A comparison analysis of Radix-2 and Radix-4 algorithmic program because it looks additional appropriate for the planning by exploitation of completely different adder architectures like RCA and CLA.

Keywords: Altered Booth encryption, Pre-Encoded multipliers, VLSI implementation.

I. INTRODUCTION

Day by day IC technology is getting more complex in terms of design and its performance analysis. A faster design with lower power consumption and smaller area is implicit to the modern electronic designs. Multipliers generally have extended latency, huge area and consume substantial amount of power. Hence low-power multiplier factor style has become a very important region in VLSI system style. Everyday new approaches are being developed to style low-power multipliers at technological, physical, circuit and logic levels. Since the multiplier is usually the slowest component in an exceedingly system, the system’s performance is decided by performance of the multiplier. Also multipliers are the most area consuming entity in a design. Therefore, optimizing speed and area of a multiplier is a major design issue nowadays. However, area and speed are usually conflicting constraints so that improving speed results in larger areas and vice-versa. Also area and power consumption of a circuit are linearly correlated. So a compromise has to be done in speed of the circuit for a greater improvement in reduction of area and power. A higher representation radix effectively indicates to fewer digits. Thus, a single-digit multiplication formula necessitates fewer cycles as we have a tendency to begin moving to a lot of higher radices that mechanically results in a lesser range of partial product. Many algorithms are developed for this purpose like Booth’s formula, Wallace Tree technique etc. For the summation method many adder architectures are accessible viz. Ripple Carry Addition, Carry Look-ahead Addition, Carry Save Addition etc. But to reduce the power consumption the summation architecture of the multiplier should be carefully chosen.

Many algorithms have been proposed for implementing efficient modular multiplication. These algorithms can be classified into the following three categories:


3. Look-up table methods: Kawamura, Takabayashi and Shimbo’s method; Hong, OH and Yoon’s method; and Lim, Hwang and Lee’s methodology.

Look-up table methods are normally faster than the generalized ones, but require a large size of memory. The Barrett algorithm and the Montgomery algorithm require small amount of pre-computation. The algorithms using pre-computation are only suitable when some parameters are fixed.

II. SYSTEM MODEL

Multiplying a variable by a group of identified constant coefficients may be a common operation in several digital signal process (DSP) algorithms. Compared to alternative common operations in DSP algorithms, like addition, subtraction, exploiting delay components, etc., multiplication is usually the foremost costly. There is a trade-off between the amount of logic resources used (i.e. the amount of silicon in the integrated circuit) and how fast the computation can be done. Compared to most of the other operations, multiplication requires more time given the same amount of logic resources and it requires more logic resources under the constraint that each operation must be completed within the same amount of time.

A general multiplier is needed if one performs multiplication between two arbitrary variables. However, when multiplying by a known constant, we can exploit the properties of binary multiplication in order to obtain a less expensive logic circuit that is functionally equivalent to simply affirming the constant on one input of a standard multiplier. In many cases, using a cheaper implementation for only multiplication still results in significant savings when considering the entire logic circuit because multiplication is relatively expensive. Furthermore, multiplication might be the vital operation, based on the application present.

A) Basic binary multiplier

The operation of multiplication is rather simple in digital electronics.

<table>
<thead>
<tr>
<th>10×8 = 80</th>
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<td>1 0 1 0</td>
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Figure 1. Basic binary multiplication algorithms.

It has its origin from the classical algorithm for the product of two binary numbers. This algorithm uses addition and shift left operations to calculate the product of two numbers. Booth's Multiplication Algorithm Booth's multiplication algorithm is a multiplication algorithm that multiplies two signed binary numbers in two's complement notation. The rule was fabricated by Andrew Donald Booth in 1950 during doing analysis on crystallography at Birkbeck faculty in Bloomsbury, London. Booth used table calculators that were quicker at shifting than adding and created the rule to extend their speed.

B) A Typical Implementation

Booth’s algorithmic program is enforced by repeatedly adding (with standard unsigned binary addition) one in every of 2 preset values A and S to a product P, then employing a rightward arithmetic shift on P. Let m and r be the number and multiplier, respectively; and let x and y represent the quantity of bits in m and r.

1) Determine the values of A and S, and the initial value of P. All of these numbers should have a length equal to (x + y + 1).
a) A: Fill the most significant (leftmost) bits with the value of m. Fill the remaining \((y + 1)\) bits with zeros.

b) S: Fill the foremost vital bits with the worth of \((-m)\) in two's complement notation. Fill the remaining \((y + 1)\) bits with zeros.

c) P: Fill the foremost vital \(x\) bits with zeros. To the proper of this, append the worth of \(r\). Fill the smallest amount vital (rightmost) bit with a zero.

2) Verify the 2 least vital (rightmost) bits of \(P\).

a) If they're 01, notice the worth of \(P + A\). Ignore any overflow.

b) If they're 10, notice the worth of \(P + S\). Ignore any overflow.

c) If they're 00, do nothing. Use \(P\) directly in the next step.

d) If they are 11, do nothing. Use \(P\) directly in the next step.

3) Arithmetically shift the value obtained in the 2nd step by a single place to the right. Let \(P\) now equal this new value.

4) Repeat steps two and three till they are been done \(y\) times.

5) Drop the smallest amount important (rightmost) bit from \(P\). This is often the product of \(m\) and \(r\).

The representations of the number and products don’t seem to be specified; usually, these are both additionally in two's complement illustration, just like the multiplier factor, however any mathematical notation that supports addition and subtraction can work still. As explicit here, the order of the steps isn't determined. Typically, it yield from LSB to MSB, beginning at \(i = 0\); the multiplication by \(2^i\) is then usually replaced by progressive shifting of the \(P\) accumulator to the right between steps; low bits are often shifted out, and consequent additions and subtractions will then be done simply on the greatest \(N\) bits.

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**III. LITERATURE REVIEW**

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<tr>
<th>Sr. No.</th>
<th>TITLE</th>
<th>AUTHOR</th>
<th>YEAR</th>
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<tr>
<td>1</td>
<td>Pre-Encoded Multipliers Based on Non-Redundant Radix-4 Signed Digit Encoding</td>
<td>K. Tsoumanis, N. Axelos, N. Moschopoulos, G. Zervakis and K. Pekmestzi,</td>
<td>2016</td>
<td>Digital signal processing applications based on off line encoding of coefficients</td>
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<td>2</td>
<td>Comparison of regular and tree based multiplier architectures with modified booth encoding for 4 bits on layout level using 45nm technology</td>
<td>B. Dinesh, V. Venkateshwaran, P. Kavinmalar and M. Kathirvelu,</td>
<td>2014</td>
<td>A detailed analysis of all the serial parallel and parallel architectures</td>
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<td>3</td>
<td>Hybrid modified booth encoded algorithm-carry save adder fast multiplier,</td>
<td>N. G. NikDaud, F. R. Hashim, M. Mustapha and M. S. Badruddin,</td>
<td>2014</td>
<td>A new architecture of hybrid Modified Booth Encoded Algorithm (MBE) and Carry Save Adder (CSA)</td>
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<td>4</td>
<td>An efficient fixed width multiplier for digital filter</td>
<td>S. Nithya and M. N. V. Nithya,</td>
<td>2014</td>
<td>A high speed and low power FIR digital filter design using the fixed width booth multiplier</td>
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<td>5</td>
<td>A New Redundant Binary Partial Product Generator for Fast 2n-Bit Multiplier Design</td>
<td>Xiaoping, H. Wei, C. Xin and W. Shumin,</td>
<td>2014</td>
<td>A new radix-16 RB Booth Encoding (RBBE-4) to avoid the hard multiple of high-radix Booth encoding without incurring any ECW.</td>
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<tr>
<td>6</td>
<td>High Speed Modified Booth Encoder Multiplier for Signed and Unsigned Numbers</td>
<td>R. P. Rajput and M. N. S. Swamy,</td>
<td>2012</td>
<td>He design and implementation of signed unsigned Modified Booth Encoding (SUMBE) multiplier.</td>
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[A] K. Tsoumanis, N. Axelos, N. Moschopoulos, G. Zervakis and K. Pekmestzi.[1] In this work, we introduce an architecture of pre-encoded multipliers for digital signal processing applications based on off-line encoding of coefficients. To this extend, the Non-Redundant radix-4 Signed-Digit (NR4SD) encoding technique, which uses the digit values $\{\text{-}1, 0, +1, +2\}$ or $\{\text{-}2, \text{-}1, 0, +1\}$, is proposed leading to a multiplier design with less complex partial products implementation. Extensive experimental analysis verifies that the planned pre-encoded NR4SD multipliers, as well as the coefficients memory, are additional region and power efficiency than the traditional altered Booth theme.

[B] Dinesh, V. Venkateshwaran, P. Kavinmalar and M. Kathirvelu,[2] Multipliers are key parts of the many high performance systems like FIR filters, microprocessors, digital signal processors, etc. A system’s performance is generally determined by the performance of the multiplier as the multiplier is generally the slowest element in the system. The analysis of performance parameters of different multiplier logics is essential for design of a system intended for a specific function with constraints on Power, Area and Delay. The work presents a detailed analysis of all the serial-parallel and parallel architectures. The multipliers are designed for 4 bit multiplication using DSCH tool and the corresponding layouts are obtained using Microwind 3.5 tool using 45nm technology. From the analysis it’s discovered that the array multipliers give a general routing structure which can be optimum for FPGA based mostly systems. Among the tree based mostly multipliers Dadda multipliers have a small advantage over Wallace tree multipliers in terms of performance. The altered booth multiplier factor is relatively inefficient for bits lesser than or up to four, because of the enhanced space concerned for realization of the booth encoder and booth selector blocks. The analysis shows that for lower order bits Dadda reduction is the most efficient.

[C] N. G. NikDaud, F. R. Hashim, M. Mustapha and M. S. Badruddin,[3] One of the effective ways to speed up multiplication are by reducing the number of partial products and accelerating the accumulation. In this work, a new architecture of hybrid Modified Booth Encoded Algorithm (MBE) and Carry Save Adder (CSA) is developed as fast multiplier architecture. Altera Quartus II platform is used to run the simulation. The architecture design is programmed into FPGA using Altera DE2 board to verify the synthesizability on physical hardware. This hybrid fast multiplier delivers good performance in term of higher speed as well as in term of less usage of logic elements.

[D] S. Nithya and M. N. V. Nithya,[4] We implement a high speed and low power FIR digital filter style employing the constant breadth booth multiplier. To scale back the miscalculation in constant breadth multiplier reconciling probability calculator is employed (ACPE). To realize higher speed, the
altered Booth encryption has been used and conjointly to fast up the addition the carry look ahead adder is employed as a carry propagate adder. The multiplier circuit is intended employing VERILOG and synthesized exploitation Xilinx ISE9.2i machine. The area, power and delay of the designed filter is analyzed exploitation cadence tool.

[E] C. Xiaoping, H. Wei, C. Xin and W. Shumin,[5] The radix-4 Booth encryption or altered Booth encryption (MBE) has been very much adopted in partial merchandise generator to style high-speed redundant binary (RB) multipliers. Attributable to the existence of the error-correcting word (ECW) generated by MBE and RB encryption, the RB multiplier generates a further RB partial product rows. An additional RB partial product accumulator (RBPAPA) stage is required for 2n-b RB MBE multiplier. The upper radix Booth formula than radix-4 is adopted to scale back the amount of partial merchandise. However, the Booth encryption isn’t economical due to the problem in generating laborious multiples. The laborious multiples drawback in RB multiplier is resolved by distinction of 2 easy power-of-two multiples. This work presents a brand new radix-16 RB Booth encryption (RBBE-4) to avoid the laborious multiple of high-radix Booth encryption while not acquisition any ECW. The planned technique results in build high-speed and low-power RB multipliers. The experimental results show that the planned RBBE-4 multiplier achieves vital improvement in delay and power consumption compared with the RB MBE multiplier and also the current accorded best RBBE-4 multipliers.

[F] R. P. Rajput and M. N. S. Swamy,[6] This work presents the planning and implementation of signed-unsigned altered Booth encryption (SUMBE) multiplier factor. This altered Booth encryption (MBE) multiplier and therefore the Baugh-Wooley multiplier performs multiplication operation on signed numbers solely. The array multiplier and Braun array multipliers perform multiplication operation on unsigned numbers solely. Thus, the necessity of the new automatic data processing system may be a dedicated and really high speed novel multiplier unit for signed and unsigned numbers. Therefore, this work presents the planning and implementation of SUMBE multiplier. The altered Booth Encoder circuit generates half the partial merchandise in parallel. By extending sign bit of the operands and generating an extra partial product the SUMBE multiplier is obtained. The Carry Save Adder (CSA) tree and therefore the final Carry Look ahead (CLA) adder accustomed speed up the multiplier operation. Since signed and unsigned multiplication operation is performed by constant multiplier unit the desired hardware and therefore the chip space reduces and this successively reduces power dissipation and price of a system.

[G] B. C. Paul, S. Fujita and M. Okajima,[7] we tend to provide a ROM-based sixteen times sixteen multiplier for low-power applications. the planning uses sixteen four times four ROM-based multiplier blocks followed by carry-save adders and a final carry-select adder (all ROM-based) to get the thirty two bit output. All memory blocks are enforced employing single semiconductor memory cells and eliminating identical rows and columns for optimizing the facility and performance. Measuring ends up in 0.18 mum CMOS method show a 400th reduction in power over the standard carry-save array multiplier once operated at its highest frequency. The ROM-based style additionally provides four hundred and forty yards less delay than the array multiplier with a minimal increase (7.7%) in power. This demonstrates the low-power operation of the ROM-based multiplier additionally at higher frequencies.

IV. PROBLEM STATEMENT

One of the numerous solutions of realizing high speed multipliers is enhancing similarity that helps in decreasing the amount of consequent calculation levels. The initial version of Booth algorithmic program (Radix-2) had 2 explicit drawbacks. They were:

International Journal of Scientific Research in Science, Engineering and Technology (ijsrset.com)
The range of add-subtract operations and shift operations become variable and causes inconvenience in realizing parallel multipliers.

The rule becomes inefficient once there are isolated 1's.

These issues are overpowered by employing altered Radix4 Booth algorithmic program that scans strings of 3 bits employing the algorithmic program given below:

1. Lengthen the sign bit 1 position if necessary to ensure that n is even.
2. Add a 0 to the right of the LSB of the multiplier. Corresponding to the value of each vector, each Partial Product will be 0, +M, -M, +2M or -2M.

V. CONCLUSION

After going through all the literature survey and review of past work and after facing a lot of problems in previous base work, we are able determine the objectives of the research work that are to implement Booth's Algorithm for the design of a binary multiplier using different architectures and power analysis at various levels. To study the space and also the time delay that consumed by distinct adders and located out an optimum relationship among the time and space complexity the adders which we have there into consideration? After examining all it might be finalized that no of bits changes are best fitted for Low Power Applications. Then the research turned focus toward the area and delay of Multipliers. Further work can be done on design a Radix-4 multiplier and estimated its delay, area.

VI. REFERENCES


